Sampling from a Convex Body

Based on

- 1. Ravi Kannan's notes.
- 2. Jonathan Kelner's lecture notes.
- 3. "Techniques in Optimization and Sampling" by Yin Tat Lee and Santosh Vempala.
- 4. "Algorithmic Convex Geometry" by Santosh Vempala.
- 5. "Geometric Random Walks" by Santosh Vempala.

Computing a point in K

Given a separation oracle for a convex body K, and $R \in \mathbb{R}_+$ such that $\mathcal{B}(y,1) \subseteq K \subseteq \mathcal{B}(0,R)$ (y is unknown), compute a point in K.

- ▶ Goal: minimize number of calls to the separation oracle.
- ► Any deterministic algorithm needs n log₂ R separation oracle calls in the worst case.

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへぐ

Computing a point in K

Given a separation oracle for a convex body K, and $R \in \mathbb{R}_+$ such that $\mathcal{B}(y,1) \subseteq K \subseteq \mathcal{B}(0,R)$ (y is unknown), compute a point in K.

- ▶ Goal: minimize number of calls to the separation oracle.
- ► Any deterministic algorithm needs n log₂ R separation oracle calls in the worst case.

Algorithm:

- 1. Initialize $P_1 = [-R, R]^n$, z = 0.
- 2. for i = 1 to N (to be fixed later.)
 - 2.1 Compute $z \in P_i$ (to be specified later).
 - 2.2 Call separation oracle on z. If $z \in K$, output z.
 - 2.3 Let $a^T x \leq b$ be the hyperplane returned by the separation oracle. Set $P_{i+1} := P_i \cap \{x \in \mathbb{R}^n : a^T x \leq a^T z\}.$
- 3. Output that K is empty.

▶ If at the end of the algorithm we have vol $(P_N) < \text{vol}(\mathcal{B}(0,1))$, then K is empty.

- If at the end of the algorithm we have vol $(P_N) < \text{vol}(\mathcal{B}(0,1))$, then K is empty.
- ▶ We would like to compute new z such that volume of P_i is guaranteed to shrink by at least a constant factor. Then $N = \Theta(n \log R)$ will suffice.

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへぐ

- If at the end of the algorithm we have vol $(P_N) < \text{vol}(\mathcal{B}(0,1))$, then K is empty.
- ▶ We would like to compute new z such that volume of P_i is guaranteed to shrink by at least a constant factor. Then $N = \Theta(n \log R)$ will suffice.

Choosing z to be the centroid guarantees this (Grunbaum's Theorem), but computing centroid is #P-hard in general.

- If at the end of the algorithm we have vol $(P_N) < \text{vol}(\mathcal{B}(0,1))$, then K is empty.
- ▶ We would like to compute new z such that volume of P_i is guaranteed to shrink by at least a constant factor. Then $N = \Theta(n \log R)$ will suffice.

Choosing z to be the centroid guarantees this (Grunbaum's Theorem), but computing centroid is #P-hard in general.

Sample *m* independent and uniform random points from P_i , denote them by y_1, \ldots, y_m . Set $z = (\sum_{i \in [m]} y_i)/m$.

- ▶ If at the end of the algorithm we have vol $(P_N) < \text{vol}(\mathcal{B}(0,1))$, then K is empty.
- We would like to compute new z such that volume of P_i is guaranteed to shrink by at least a constant factor. Then $N = \Theta(n \log R)$ will suffice.
- Choosing z to be the centroid guarantees this (Grunbaum's Theorem), but computing centroid is #P-hard in general.

Sample *m* independent and uniform random points from P_i , denote them by y_1, \ldots, y_m . Set $z = (\sum_{i \in [m]} y_i)/m$.

Theorem

$$\mathbb{E}\left[ext{vol}\left(P_{i+1}
ight)
ight] \leq \left(1 - rac{1}{e} + \sqrt{rac{n}{m}}
ight) ext{vol}\left(P_{i}
ight).$$

With $m = \Theta(n)$, $N = \Theta(n \log R)$ will suffice.

For a set of vertices $S \subset V$,

$$\phi(S) := \frac{\sum_{i \in S, j \in V \setminus S} \pi_i P_{ij}}{\min \{\pi(S), \pi(V \setminus S)\}}$$

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへで

and $\phi := \min_{S:0 < \pi(S) < 1} \phi(S)$.

For a set of vertices $S \subset V$,

$$\phi(S) := \frac{\sum_{i \in S, j \in V \setminus S} \pi_i P_{ij}}{\min \left\{ \pi(S), \pi(V \setminus S) \right\}}$$

and $\phi := \min_{S:0 < \pi(S) < 1} \phi(S)$. Markov Chain (informal definition)

► (K, A) where K is the state space and A is a set of subsets of K that is closed under complements and countable unions.

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへぐ

For a set of vertices $S \subset V$,

$$\phi(S) := \frac{\sum_{i \in S, j \in V \setminus S} \pi_i P_{ij}}{\min \left\{ \pi(S), \pi(V \setminus S) \right\}}$$

and $\phi := \min_{S:0 < \pi(S) < 1} \phi(S)$. Markov Chain (informal definition)

- ► (K, A) where K is the state space and A is a set of subsets of K that is closed under complements and countable unions.
- For each u ∈ K and A ∈ A, P_u(A) is the probability of being in A after taking one step from u.

For a set of vertices $S \subset V$,

$$\phi(S) := \frac{\sum_{i \in S, j \in V \setminus S} \pi_i P_{ij}}{\min \left\{ \pi(S), \pi(V \setminus S) \right\}}$$

and $\phi := \min_{S:0 < \pi(S) < 1} \phi(S)$. Markov Chain (informal definition)

- ► (K, A) where K is the state space and A is a set of subsets of K that is closed under complements and countable unions.
- For each u ∈ K and A ∈ A, P_u(A) is the probability of being in A after taking one step from u.
- Given a starting distribution Q_0 , w_0 is sampled from Q_0 and w_i is sampled from $P_{w_{i-1}}$.

A distribution Q on (K, A) is called stationary if one step from it gives the same distribution, i.e., for any A ∈ A,

$$\int_{K} P_u(A) dQ(u) = Q(A).$$

A distribution Q on (K, A) is called stationary if one step from it gives the same distribution, i.e., for any A ∈ A,

$$\int_{K} P_u(A) dQ(u) = Q(A).$$

• The conductance of a subset A is defined as

$$\phi(A) := \frac{\int_A P_u(K \setminus A) dQ(u)}{\min \{Q(A), Q(K \setminus A)\}}$$

and the conductance of the Markov chain is $\phi := \min_A \phi(A)$.

◆□▶ ◆□▶ ◆ □▶ ◆ □ ● ● ● ●

A distribution Q on (K, A) is called stationary if one step from it gives the same distribution, i.e., for any A ∈ A,

$$\int_{\mathcal{K}} P_u(A) dQ(u) = Q(A).$$

► The conductance of a subset A is defined as

$$\phi(A) := \frac{\int_A P_u(K \setminus A) dQ(u)}{\min \{Q(A), Q(K \setminus A)\}}$$

and the conductance of the Markov chain is $\phi := \min_A \phi(A)$.

- A distribution Q is *atom-free* if there is no $x \in K$ with Q(x) > 0.
- ▶ *P* is said to be *M*-warm with respect to *Q* if

$$M = \sup_{A \in \mathcal{A}} \frac{P(A)}{Q(A)}.$$

Mixing Time

Let Q_t be the distribution of the random walk at time t.

Theorem Let $M = \sup_A Q_0(A)/Q(A)$. Then

$$d_{\mathsf{TV}}\left(\mathcal{Q}_t, \mathcal{Q}
ight) \leq \sqrt{M} \left(1 - rac{\phi^2}{2}
ight)^t.$$

Mixing Time

Let Q_t be the distribution of the random walk at time t.

Theorem Let $M = \sup_A Q_0(A)/Q(A)$. Then

$$d_{\mathsf{TV}}\left(\mathcal{Q}_t, \mathcal{Q}
ight) \leq \sqrt{\mathcal{M}}\left(1 - rac{\phi^2}{2}
ight)^t.$$

Therefore, for

$$t = O\left(rac{1}{\phi^2}\log\left(rac{M}{\epsilon}
ight)
ight) \qquad ext{we have} \qquad d_{\mathsf{TV}}\left(Q_t,Q
ight) \leq \epsilon.$$

Mixing Time

Let Q_t be the distribution of the random walk at time t.

Theorem Let $M = \sup_A Q_0(A)/Q(A)$. Then

$$d_{\mathsf{TV}}\left(Q_t, Q
ight) \leq \sqrt{M} \left(1 - rac{\phi^2}{2}
ight)^t.$$

Therefore, for

$$t = O\left(rac{1}{\phi^2}\log\left(rac{M}{\epsilon}
ight)
ight) \qquad ext{we have} \qquad d_{\mathsf{TV}}\left(\mathcal{Q}_t, \mathcal{Q}
ight) \leq \epsilon.$$

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへで

To bound mixing time, sufficient to prove lower bound on ϕ .

Fix a set $S \subset K$ and consider the Grid-walk or ball-walk where the walk is currently in S.

Fix a set $S \subset K$ and consider the Grid-walk or ball-walk where the walk is currently in S. In one step, walk has a "large" probability of going to $K \setminus S$ if walk is currently "close" to the boundary of S and $K \setminus S$.

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへぐ

Fix a set $S \subset K$ and consider the Grid-walk or ball-walk where the walk is currently in S. In one step, walk has a "large" probability of going to $K \setminus S$ if walk is currently "close" to the boundary of S and $K \setminus S$.

For $S_1, S_2 \subseteq K$,

$$d(S_1, S_2) := \inf \{ \|u - v\| : u \in S_1, \ v \in S_2 \}.$$

Fix a set $S \subset K$ and consider the Grid-walk or ball-walk where the walk is currently in S. In one step, walk has a "large" probability of going to $K \setminus S$ if walk is currently "close" to the boundary of S and $K \setminus S$.

For
$$S_1, S_2 \subseteq K$$
, $d(S_1, S_2) := \inf \{ \|u - v\| : u \in S_1, v \in S_2 \}$.

Theorem

Let S_1, S_2, S_3 be a partition into measurable sets of a convex body K of diameter D. Then,

$$\operatorname{vol}(S_3) \geq rac{2d(S_1, S_2)}{D} \min \left\{ \operatorname{vol}(S_1), \operatorname{vol}(S_2) \right\}.$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ の 0 0

Fix a set $S \subset K$ and consider the Grid-walk or ball-walk where the walk is currently in S. In one step, walk has a "large" probability of going to $K \setminus S$ if walk is currently "close" to the boundary of S and $K \setminus S$.

For
$$S_1, S_2 \subseteq K$$
, $d(S_1, S_2) := \inf \{ \|u - v\| : u \in S_1, v \in S_2 \}$.

Theorem

Let S_1, S_2, S_3 be a partition into measurable sets of a convex body K of diameter D. Then,

$$\operatorname{vol}(S_3) \geq rac{2d(S_1, S_2)}{D} \min \left\{ \operatorname{vol}(S_1), \operatorname{vol}(S_2) \right\}.$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ の 0 0

▶ Not true for non-convex bodies (e.g. dumbell graph).

Fix a set $S \subset K$ and consider the Grid-walk or ball-walk where the walk is currently in S. In one step, walk has a "large" probability of going to $K \setminus S$ if walk is currently "close" to the boundary of S and $K \setminus S$.

For
$$S_1, S_2 \subseteq K$$
,
 $d(S_1, S_2) := \inf \{ \|u - v\| : u \in S_1, v \in S_2 \}.$

Theorem

Let S_1, S_2, S_3 be a partition into measurable sets of a convex body K of diameter D. Then,

$$\operatorname{\mathsf{vol}}\left(S_3
ight) \geq rac{2d(S_1,S_2)}{D}\min\left\{\operatorname{\mathsf{vol}}\left(S_1
ight),\operatorname{\mathsf{vol}}\left(S_2
ight)
ight\}.$$

- ▶ Not true for non-convex bodies (e.g. dumbell graph).
- ► Dependance on *D* unavoidable (e.g. cylinder verify this).

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be an integrable function. It defines a measure π_f on measurable subsets of \mathbb{R}^n .

$$\pi_f(A) = \frac{\int_A f(x) dx}{\int_{\mathbb{R}^n} f(x) dx}.$$

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be an integrable function. It defines a measure π_f on measurable subsets of \mathbb{R}^n .

$$\pi_f(A) = \frac{\int_A f(x) dx}{\int_{\mathbb{R}^n} f(x) dx}$$

Ball walk with Metropolis filter (δ, f)

1. Pick a uniform random point in the ball of radius δ centered at the current point x.

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへで

2. Move to y with probability min $\{1, f(y)/f(x)\}$, stay at x with remaining probability.

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be an integrable function. It defines a measure π_f on measurable subsets of \mathbb{R}^n .

$$\pi_f(A) = \frac{\int_A f(x) dx}{\int_{\mathbb{R}^n} f(x) dx}$$

Ball walk with Metropolis filter (δ, f)

1. Pick a uniform random point in the ball of radius δ centered at the current point x.

- 2. Move to y with probability min $\{1, f(y)/f(x)\}$, stay at x with remaining probability.
- Ball walk has π_f as stationary distribution.

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be an integrable function. It defines a measure π_f on measurable subsets of \mathbb{R}^n .

$$\pi_f(A) = \frac{\int_A f(x) dx}{\int_{\mathbb{R}^n} f(x) dx}$$

Ball walk with Metropolis filter (δ, f)

- 1. Pick a uniform random point in the ball of radius δ centered at the current point x.
- 2. Move to y with probability min $\{1, f(y)/f(x)\}$, stay at x with remaining probability.
- Ball walk has π_f as stationary distribution.
- ▶ When does this have "large" conductance?

Logconcave functions

A function $f : \mathbb{R}^n \to \mathbb{R}_+$ is said to be *logconcave* if for any $x, y \in \mathbb{R}^n$ and $\lambda \in [0, 1]$,

 $f(\lambda x + (1-\lambda)y) \ge f(x)^{\lambda}f(y)^{1-\lambda}.$

Logconcave functions

A function $f : \mathbb{R}^n \to \mathbb{R}_+$ is said to be *logconcave* if for any $x, y \in \mathbb{R}^n$ and $\lambda \in [0, 1]$,

$$f(\lambda x + (1-\lambda)y) \ge f(x)^{\lambda}f(y)^{1-\lambda}.$$

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ ● のへで

Examples (verify)

1. For a convex body K, let
$$f(x) = \begin{cases} 1 & \text{if } x \in K \\ 0 & \text{otherwise} \end{cases}$$

- 2. $f(x) = e^{-||x||^2}$.
- 3. Product of two logconcave functions.

Theorem

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be a logconcave density function, and let H be any halfspace containing its centroid. Then

$$\int_{H} f(x) dx \geq \frac{1}{e}.$$

Theorem

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be a logconcave density function, and let H be any halfspace containing its centroid. Then

$$\int_{H} f(x) dx \geq \frac{1}{e}$$

Theorem

Let $f : \mathbb{R}^n \to \mathbb{R}_+$ be a logconcave density function, z be the average of m independent random points from π_f , and let H be any halfspace containing z. Then

$$\mathbb{E}\left[\pi_f(H)\right] \geq \frac{1}{e} - \sqrt{\frac{n}{m}}.$$

Theorem

Let f be a logconcave density on \mathbb{R}^n whose support has diameter D. Then for any partition of \mathbb{R}^n into measurable sets S_1, S_2, S_3

$$\pi_f(S_3) \geq \frac{2d(S_1, S_2)}{D} \min \{\pi_f(S_1), \pi_f(S_2)\}.$$

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへで

Theorem

Let f be a logconcave density on \mathbb{R}^n whose support has diameter D. Then for any partition of \mathbb{R}^n into measurable sets S_1, S_2, S_3

$$\pi_f(S_3) \geq \frac{2d(S_1, S_2)}{D} \min \{\pi_f(S_1), \pi_f(S_2)\}.$$

Let z_f be the centroid of f and let $M(f) := \mathbb{E}_{x \sim \pi_f} \|x - z_f\|$.

Theorem

Let f be a logconcave density on \mathbb{R}^n whose support has diameter D. Then for any partition of \mathbb{R}^n into measurable sets S_1, S_2, S_3

$$\pi_f(S_3) \geq \frac{2d(S_1, S_2)}{D} \min \{\pi_f(S_1), \pi_f(S_2)\}.$$

Let z_f be the centroid of f and let $M(f) := \mathbb{E}_{x \sim \pi_f} \|x - z_f\|$.

Theorem

Let f be a logconcave density on \mathbb{R}^n . Then for any partition of \mathbb{R}^n into measurable sets S_1, S_2, S_3

$$\pi_f(S_3) \geq \frac{\ln 2}{M(f)} d(S_1, S_2) \min \{\pi_f(S_1), \pi_f(S_2)\}.$$

Isotropic Densities

A density function $f : \mathbb{R}^n \to \mathbb{R}_+$ is *isotropic* if

$$\mathbb{E}_{X \sim \pi_f} [X] = 0$$
 and $\mathbb{E}_{X \sim \pi_f} [XX^T] = I.$

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへで

Isotropic Densities

A density function $f : \mathbb{R}^n \to \mathbb{R}_+$ is *isotropic* if

$$\mathbb{E}_{X \sim \pi_f} [X] = 0$$
 and $\mathbb{E}_{X \sim \pi_f} [XX^T] = I.$

For an isotropic density f,

$$M(f) = \mathbb{E}_{X \sim \pi_f} \|X\| \leq \sqrt{\mathbb{E}_{X \sim \pi_f} \|X\|^2} = \sqrt{n}.$$

Isotropic Densities

A density function $f : \mathbb{R}^n \to \mathbb{R}_+$ is *isotropic* if

$$\mathbb{E}_{X \sim \pi_f} [X] = 0$$
 and $\mathbb{E}_{X \sim \pi_f} [XX^T] = I.$

For an isotropic density f,

$$M(f) = \mathbb{E}_{\boldsymbol{X} \sim \pi_f} \left\| \boldsymbol{X}
ight\| \leq \sqrt{\mathbb{E}_{\boldsymbol{X} \sim \pi_f} \left\| \boldsymbol{X}
ight\|^2} = \sqrt{n}.$$

Theorem For an isotropic logconcave density f,

 $\mathbb{P}\left[\|X\| \geq t\sqrt{n}\right] \leq e^{-t}.$

▲□▶ ▲□▶ ▲□▶ ▲□▶ □ の 0 0

$$\mathbb{E}_{X \sim \pi_f} \left[X \right] = 0 \qquad \text{and} \qquad \frac{1}{C} I \preceq \mathbb{E}_{X \sim \pi_f} \left[X X^T \right] \preceq C I.$$

$$\mathbb{E}_{X \sim \pi_f} [X] = 0$$
 and $\frac{1}{C}I \preceq \mathbb{E}_{X \sim \pi_f} [XX^T] \preceq CI.$

• Let
$$\Sigma = \mathbb{E}_{X \sim \pi_f} \left[(X - \mathbb{E}[X]) (X - \mathbb{E}[X])^T \right]$$
. Then $Y := \Sigma^{-1/2} (X - \mathbb{E}[X])$ is isotropic.

$$\mathbb{E}\left[YY^{T}\right] = \mathbb{E}\left[\Sigma^{-1/2}\left(X - \mathbb{E}\left[X\right]\right)\left(\Sigma^{-1/2}\left(X - \mathbb{E}\left[X\right]\right)\right)^{T}\right]$$
$$= \Sigma^{-1/2}\mathbb{E}\left[\left(X - \mathbb{E}\left[X\right]\right)\left(X - \mathbb{E}\left[X\right]\right)^{T}\right]\Sigma^{-1/2} = I.$$

$$\mathbb{E}_{X \sim \pi_f} \left[X \right] = 0 \qquad ext{and} \qquad rac{1}{C} I \preceq \mathbb{E}_{X \sim \pi_f} \left[X X^T \right] \preceq C I.$$

• Let
$$\Sigma = \mathbb{E}_{X \sim \pi_f} \left[(X - \mathbb{E}[X]) (X - \mathbb{E}[X])^T \right]$$
. Then $Y := \Sigma^{-1/2} (X - \mathbb{E}[X])$ is isotropic.

$$\mathbb{E}\left[YY^{T}\right] = \mathbb{E}\left[\Sigma^{-1/2}\left(X - \mathbb{E}\left[X\right]\right)\left(\Sigma^{-1/2}\left(X - \mathbb{E}\left[X\right]\right)\right)^{T}\right]$$
$$= \Sigma^{-1/2}\mathbb{E}\left[\left(X - \mathbb{E}\left[X\right]\right)\left(X - \mathbb{E}\left[X\right]\right)^{T}\right]\Sigma^{-1/2} = I.$$

▶ For any linear transformation A, $vol(AK) = det(A) \cdot vol(K)$.

$$\mathbb{E}_{X \sim \pi_f} \left[X \right] = 0 \qquad ext{and} \qquad rac{1}{C} I \preceq \mathbb{E}_{X \sim \pi_f} \left[X X^T \right] \preceq C I.$$

• Let
$$\Sigma = \mathbb{E}_{X \sim \pi_f} \left[(X - \mathbb{E}[X]) (X - \mathbb{E}[X])^T \right]$$
. Then $Y := \Sigma^{-1/2} (X - \mathbb{E}[X])$ is isotropic.

$$\mathbb{E}\left[YY^{T}\right] = \mathbb{E}\left[\Sigma^{-1/2}\left(X - \mathbb{E}\left[X\right]\right)\left(\Sigma^{-1/2}\left(X - \mathbb{E}\left[X\right]\right)\right)^{T}\right]$$
$$= \Sigma^{-1/2}\mathbb{E}\left[\left(X - \mathbb{E}\left[X\right]\right)\left(X - \mathbb{E}\left[X\right]\right)^{T}\right]\Sigma^{-1/2} = I.$$

▲□▶ ▲□▶ ▲ □▶ ▲ □▶ □ のへで

- ▶ For any linear transformation A, $vol(AK) = det(A) \cdot vol(K)$.
- Σ can be estimated by sampling?