

What if the convex function is Lipschitz?

Assumptions: ① $\|\nabla f(x)\|_2 \leq G \forall x$, ② $\|x^0 - x^*\|_2 \leq D$

Theorem: Grad descent with an appropriate choice of step-size η ,

will in $T = O\left(\frac{GD}{\epsilon}\right)^2$ iterations output a sequence of points x^0, x^1, \dots, x^{T-1} s.t.

$$f\left(\frac{1}{T} \sum_{t=0}^{T-1} x^t\right) - f(x^*) \leq \epsilon$$

Remarks ① Different convergence guarantee than the smooth setting (where we had $f(x^T) - f(x^*) \leq \epsilon$).

② worse convergence guarantee ($\frac{1}{\epsilon}$ vs $\frac{1}{\epsilon^2}$). This is tight ($\Omega\left(\frac{GD}{\epsilon}\right)^2$ lower bound in the first ^{order} oracle setting).

→ What if $\|\nabla f(x)\|_\infty$ was bounded (as opposed to $\|\nabla f(x)\|_2$)?

A very general recipe/algorithm: mirror descent that studies different bounded norm constraints on the gradient.

In this lecture, will study a special case: exponential gradient descent.

$$K = \Delta_n := \left\{ p \in [0,1]^n : \sum_{i=1}^n p_i = 1 \right\} \quad (\text{simplex})$$

$$f: \Delta_n \rightarrow \mathbb{R} \quad \text{convex, } \min_{p \in \Delta_n} f(p)$$

A general strategy for optimization:

Construct a sequence of approximations f_t to f .

$$\& \text{ set } p^{t+1} := \underset{p \in \Delta_n}{\operatorname{argmin}} f_t(p)$$

e.g. the gradient descent for smooth convex functions:

$$f_t(x) = f(x_t) + \langle \nabla f(x_t), x - x_t \rangle + \frac{L}{2} \|x - x_t\|^2$$
$$f(x) \leq f_t(x) \quad (\text{if } f \text{ is } L\text{-smooth})$$

$$f(p_t) + \langle \nabla f(p_t), p - p_t \rangle \leq f(p) \quad (\text{by convexity of } f)$$

$f_t(p)$?

What about $p_{t+1} := \underset{p \in \Delta_n}{\operatorname{argmin}} (f(p_t) + \langle \nabla f(p_t), p - p_t \rangle)$?

Too aggressive: keeps oscillating between the vertices.
So regularize!

$$p_{t+1} := \underset{p \in \Delta_n}{\operatorname{argmin}} (M(p, p_t) + f(p_t) + \langle \nabla f(p_t), p - p_t \rangle)$$

M is some sort of distance measure on the simplex.

KL divergence: Let $p, q \in \Delta_n$. Then

$$D(p \parallel q) = \sum_i p_i \log \left(\frac{p_i}{q_i} \right)$$

(lets say
log is
the natural
log)

Nice properties: ① $D(p||q) \geq 0$ the natural log
 ② $D(p||q) \geq \frac{1}{2} \|p - q\|_1^2$ (Pinsker's ineq.)

③ Jointly convex in p, q

$$D((1-\lambda)p_1 + \lambda p_2 || (1-\lambda)q_1 + \lambda q_2) \leq (1-\lambda) D(p_1 || q_1) + \lambda D(p_2 || q_2)$$

$$b_{t+1} := \operatorname{argmin}_{b \in \Delta_n} \left\{ D(b || p_t) + \eta (f(p_t) + \langle \nabla f(p_t), b - p_t \rangle) \right\}$$

$$= \operatorname{argmin}_{b \in \Delta_n} \left\{ D(b || p_t) + \eta \langle \nabla f(p_t), b \rangle \right\}$$

Lemma: $q \in \Delta_n, g \in \mathbb{R}^n$. Let $b^* = \operatorname{argmin}_{b \in \Delta_n} \left\{ D(b || q) + \eta \langle g, b \rangle \right\}$

Then $b^* = \frac{w^*}{\|w^*\|_1}$, where $w_i^* = q_i e^{-\eta g_i}$

Proof: $D(b || q) + \eta \langle g, b \rangle - D(b^* || q) - \eta \langle g, b^* \rangle$

$$= D(p || q) + \eta \langle g, b \rangle - \sum_i b_i^* \log \left(\frac{b_i^*}{q_i} \right) - \eta \langle g, b^* \rangle$$

$$= D(p || q) + \eta \langle g, b \rangle - \sum_i b_i^* \log \left(\frac{e^{-\eta g_i}}{\|w^*\|_1} \right) - \eta \langle g, b^* \rangle$$

$$= D(p || q) + \eta \langle g, b \rangle + \log(\|w^*\|_1)$$

∴ ∴

$$\begin{aligned}
 D(p \| p^*) &= \sum_i p_i \log \left(\frac{p_i}{p_i^*} \right) \\
 &= \sum_i p_i \log \left(\frac{p_i \|w^*\|_1}{z_i e^{-\eta g_i}} \right) \\
 &= D(p \| q) + \log(\|w^*\|_1) + \eta \langle g, p \rangle
 \end{aligned}$$

Hence $D(p \| q) + \eta \langle g, p \rangle - D(p^* \| q) - \eta \langle g, p^* \rangle$

$$= D(p \| p^*) \geq 0$$

↳ Jensen's inequality & concavity of log

Algorithm (Exponential gradient descent):

$$p^0 := \frac{1}{n} \mathbf{1} \text{ (uniform distribution)}$$

for $t = 0, 1, \dots, T-1$

$$g^t := \nabla f(p^t)$$

$$w_i^{t+1} := p_i^t e^{-\eta g_i^t}$$

$$p_i^{t+1} := \frac{w_i^{t+1}}{\|w^{t+1}\|_1}$$

return $\bar{p} = \frac{1}{T} \sum_{t=0}^{T-1} p^t$

Assumption: $\|\nabla f(p)\|_\infty \leq G \forall p \in \Delta_n$

Thm: with an appropriate choice of η & $T = \Theta\left(\frac{G^2 \log(n)}{\epsilon^2}\right)$

EGD returns \bar{p} s.t.

$$f(\bar{p}) - f(p^*) \leq \epsilon$$

$$(p^* = \operatorname{argmin}_{p \in \Delta_n} f(p))$$

Jensen's inequality

Proof:

$$f(\bar{p}) - f(p^*) \stackrel{\text{Jensen's Ineq.}}{\leq} \left(\frac{1}{T} \sum_{t=0}^{T-1} f(p^t) \right) - f(p^*)$$

$$= \frac{1}{T} \sum_{t=0}^{T-1} (f(p^t) - f(p^*))$$

Convexity of f

$$\leq \frac{1}{T} \sum_{t=0}^{T-1} \langle \nabla f(p^t), p^t - p^* \rangle$$

$$= \frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p^t - p^* \rangle$$

Strategy: measure decrease in the KL divergence to the optimal solution

$$D(p^* \| p^t) - D(p^* \| p^{t+1})$$

$$= D(p^* \| p^t) - \sum_i p_i^* \log \left(\frac{p_i^*}{p_i^{t+1}} \right)$$

$$= D(p^* \| p^t) - \sum_i p_i^* \log \left(\frac{p_i^*}{p_i^t e^{-\eta g_i^t}} \right)$$

$$= \eta \langle g^t, p^* \rangle - \log \left(\|w^{t+1}\|_1 \right)$$

$$\|w^{t+1}\|_1 = \sum_i p_i^t e^{-\eta g_i^t}$$

$$\log \left(\sum_i p_i^t e^{-\eta g_i^t} \right) \stackrel{\text{Jensen's}}{\geq} \sum_i p_i^t \log \left(e^{-\eta g_i^t} \right) = -\eta \langle g^t, p^t \rangle$$

↳ want to lower bound the decrease in KL divergence

& hence need an upper bound on this term

Useful fact: $e^{-x} \leq 1 - x + \frac{ex^2}{2} \quad \forall |x| \leq 1$

We know: $\|g^t\|_\infty \leq G$

$$\Rightarrow e^{-\eta g_i^t} \leq 1 - \eta g_i^t + \frac{e\eta^2 G^2}{2} \text{ if } \eta \leq \frac{1}{G}.$$

(using the above fact)

$$\Rightarrow \sum_i p_i^t e^{-\eta g_i^t} \leq 1 - \eta \langle g^t, p^t \rangle + \frac{e\eta^2 G^2}{2}$$

$$\Rightarrow \log\left(\sum_i p_i^t e^{-\eta g_i^t}\right) \leq -\eta \langle g^t, p^t \rangle + \frac{e\eta^2 G^2}{2}$$

($\log(1+y) \leq y$)

Hence $D(p^* \| p^t) - D(p^* \| p^{t+1})$
 $\geq \eta \langle g^t, p^t - p^* \rangle - \frac{e\eta^2 G^2}{2}$

summing up & telescoping

$$D(p^* \| p^0) - D(p^* \| p^T) \geq \eta \sum_{t=0}^{T-1} \langle g^t, p^t - p^* \rangle - \frac{e\eta^2 G^2 T}{2}$$

$$\frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p^t - p^* \rangle \leq \frac{D(p^* \| p^0) - D(p^* \| p^T)}{\eta T} + \frac{e\eta G^2}{2}$$

$\log(n) \leftarrow \dots \rightarrow \| \cdot \|_1$

$$\leq \frac{D(p^* \| p^0)}{\eta T} + \frac{\epsilon \eta G^2}{2}$$

$$D(p^* \| p^0) = \log(n) - H(p^*) \leq \log(n)$$

$$\begin{aligned} \hookrightarrow \sum_i p_i^* \log\left(\frac{p_i^*}{1/n}\right) &= \log(n) + \sum_i p_i^* \log(p_i^*) \\ &= \log(n) - \sum_i p_i^* \log\left(\frac{1}{p_i^*}\right) \geq 0 \end{aligned}$$

$$\leq \frac{\log(n)}{\eta T} + \frac{\epsilon \eta G^2}{2}$$

pick η so that $\frac{\log(n)}{\eta T} = \frac{\epsilon \eta G^2}{2} \Rightarrow \eta = \Theta\left(\frac{\sqrt{\log(n)}}{G \sqrt{T}}\right)$

plugging this η

$$\Theta\left(\frac{G \sqrt{\log(n)}}{\sqrt{T}}\right) \rightarrow \text{we want this} \leq \epsilon$$

so pick $T = \Theta\left(\frac{G^2 \log(n)}{\epsilon^2}\right)$

This completes the proof.

Actually can replace p^* with any p & g^t 's need not be gradients of some f .

Algorithm: Multiplicative weights update (MWU)

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essentially the same as EGD

Just that some oracle is providing these g^t 's.

$$\frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p^t \rangle - \operatorname{argmin}_{p \in \Delta_n} \left(\frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p \rangle \right) \leq \varepsilon$$

Assumption: $\|g^t\|_\infty \leq G$.

Example setting:

- n experts in the stock market.
- Not sure which expert to follow

On day t , expert i incurs a loss of g_i^t .

Use MWU to predict their credibility from previous days data.

p^t only depends on g^1, \dots, g^{t-1}

on day t , invest probabilistically acc. to p^t .

total expected loss (average over the days): $\frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p^t \rangle$

Guarantee: $\frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p^t \rangle - \operatorname{argmin}_{p \in \Delta_n} \frac{1}{T} \sum_{t=0}^{T-1} \langle g^t, p \rangle \leq \varepsilon$

regret